Windows File Access (1)

- On a low level Windows provides the same file access functions as Linux / Unix, and they have the same syntax
 - fd=open (), close (fd)
 - read (fd,...), write (fd,...)
- fopen, fread exist on Windows as well
- however, the standard way is to use Windowsspecific file functions

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9. File Systems (4) - Slide 3

Windows File Access (2)

- Windows standard functions are object-based:
 - handle = CreateFile () creates / opens a file
 - **ReadFile** (handle, ...) reads from a file
 - WriteFile (handle, ...) writes into a file
 - CloseFile (handle) closes a file
 - SetFilePointer (handle,...)
 change read/write position
 - WaitForSingleObject (handle)
 wait for termination of a file I/O operation

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Windows File Access (3)

• open / create file: CreateFile

```
HANDLE CreateFile(
   LPCTSTR lpFileName,
   DWORD dwDesiredAccess,
   DWORD dwShareMode,
   LPSECURITY_ATTRIBUTES lpSecurityAttributes,
   DWORD dwCreationDisposition,
   DWORD dwFlagsAndAttributes,
   HANDLE hTemplateFile);
```

example 1: open a file for reading

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9. File Systems (4) - Slide 5

Windows File Access (4)

example 2: open a file in append mode

append data of ONE.TXT to TWO.TXT; with locking

```
do {
   if (ReadFile(hFile, buff, 4096, &dwBytesRead, NULL)) {
      dwPos = SetFilePointer(hAppend, 0, NULL, FILE_END);
      LockFile(hAppend, dwPos, 0, dwPos + dwBytesRead, 0);
      WriteFile(hAppend, buff, dwBytesRead, &dwBytesWritten, NULL);
      UnlockFile(hAppend, dwPos, 0, dwPos + dwBytesRead, 0);
   }
} while (dwBytesRead == 4096);
// Close both files.
CloseHandle(hFile);
CloseHandle(hAppend);
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9. File Systems (4) - Slide 6
```

Windows File Access (5)

reading: ReadFile

```
BOOL ReadFile(
HANDLE hFile,
LPVOID lpBuffer,
DWORD nNumberOfBytesToRead,
LPDWORD lpNumberOfBytesRead,
LPOVERLAPPED lpOverlapped);
```

writing: WriteFile

```
BOOL WriteFile(
HANDLE hFile,
LPCVOID lpBuffer,
DWORD nNumberOfBytesToWrite,
LPDWORD lpNumberOfBytesWritten,
LPOVERLAPPED lpOverlapped);
```

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9. File Systems (4) - Slide 7

Windows File Access (6)

example: copy a file and convert the data

```
#include <windows.h>
#include <stdio.h>
#define BUFSIZE 512
 HANDLE hFile, hTempFile;
 DWORD dwRetVal, dwBytesRead, dwBytesWritten, dwBufSize=BUFSIZE;
 UINT uRetVal; BOOL fSuccess;
 char szTempName[BUFSIZE], buffer[BUFSIZE], lpPathBuffer[BUFSIZE];
 // Open the existing file.
 hFile = CreateFile("original.txt",
                                            // file name
                    GENERIC READ.
                                            // open for reading
                                            // do not share
                                            // default security
                    OPEN EXISTING,
                                            // existing file only
                    FILE_ATTRIBUTE_NORMAL, // normal file
                    NULL);
                                            // no template
 if (hFile == INVALID_HANDLE_VALUE) {
     printf ("CreateFile failed with error %d.\n", GetLastError()); return (1);
```

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Windows File Access (7)

```
// Get the temp path.
dwRetVal = GetTempPath(dwBufSize,
                                    // length of the buffer
                      lpPathBuffer); // buffer for path
if (dwRetVal > dwBufSize | | (dwRetVal == 0)) {
   printf ("GetTempPath failed with error %d.\n", GetLastError()); return (2);
// Create a temporary file.
uRetVal = GetTempFileName(lpPathBuffer, // directory for tmp files
                          "NEW".
                                       // temp file name prefix
                                       // create unique name
                          szTempName); // buffer for name
if (uRetVal == 0) {
    printf ("GetTempFileName failed with error %d.\n", GetLastError()); return (3);
// Create the new file to write the upper-case version to.
hTempFile = CreateFile((LPTSTR) szTempName, // file name
                       GENERIC_READ | GENERIC_WRITE, // open r-w
                                            // do not share
                                             // default security
                       CREATE ALWAYS,
                                            // overwrite existing
                       FILE ATTRIBUTE_NORMAL,// normal file
                       NULL);
if (hTempFile == INVALID HANDLE VALUE) {
    printf ("CreateFile failed with error %d.\n", GetLastError()); return (4);
```

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9. File Systems (4) – Slide 9

Windows File Access (8)

```
// Read BUFSIZE blocks to the buffer. Change all characters in the buffer to
       // upper case. Write the buffer to the temporary file.
       do ·
              (ReadFile(hFile, buffer, BUFSIZE, &dwBytesRead, NULL)) {
                CharUpperBuff(buffer, dwBytesRead);
               fSuccess = WriteFile(hTempFile, buffer, dwBytesRead,
                                     &dwBvtesWritten, NULL);
               if (!fSuccess)
                   printf ("WriteFile failed with error %d.\n", GetLastError()); return (5);
           } else {
               printf ("ReadFile failed with error %d.\n", GetLastError()); return (6);
       } while (dwBytesRead == BUFSIZE);
       // Close the handles to the files.
       fSuccess = CloseHandle (hFile);
       if (!fSuccess) {
          printf ("CloseHandle failed with error %d.\n", GetLastError()); return (7);
       fSuccess = CloseHandle (hTempFile);
          printf ("CloseHandle failed with error %d.\n", GetLastError()); return (8);
       // Move the temporary file to the new text file.
       fSuccess = MoveFileEx(szTempName, "allcaps.txt", MOVEFILE_REPLACE_EXISTING);
           printf ("MoveFileEx failed with error %d.\n", GetLastError()); return (9);
       return (0);
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                                                                         9. File Systems (4) - Slide 10
```

Windows File Access (9)

• jump to a specific position within the file

```
DWORD SetFilePointer(
HANDLE hFile,
LONG lDistanceToMove,
PLONG lpDistanceToMoveHigh,
DWORD dwMoveMethod
);
```

hFile: [in] A handle to the file that has a file pointer to be moved.

IDistanceToMove: [in] The low order 32-bits of a signed value that specifies the number of bytes to move the file pointer.

IpDistanceToMoveHigh: [in, out, optional] A pointer to the high order 32-bits of the signed 64-bit distance to move. If you do not need the high order 32-bits, this pointer must be set to NULL.

dwMoveMethod: [in] The starting point for the file pointer move:

FILE_BEGIN
FILE_CURRENT
FILE_END
The starting point is 0 (zero) or the beginning of the file.
The starting point is the current value of the file pointer.
The starting point is the current end-of-file position.

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9. File Systems (4) - Slide 11

Searching for Files (1)

hFind = FindFirstFile ()

```
HANDLE FindFirstFile(
   LPCTSTR lpFileName,
   LPWIN32_FIND_DATA lpFindFileData
);
```

IpFileName: [in] A pointer to a null-terminated string that specifies a valid directory or path, and file name that can contain wildcard characters, for example, an asterisk (*) or a question mark (?).

IpFindFileData: [out] A pointer to the WIN32_FIND_DATA structure that receives information about a found file or subdirectory.

FindNextFile (hFind)

```
BOOL FindNextFile(
HANDLE hFindFile,
LPWIN32_FIND_DATA lpFindFileData
```

FindClose (hFind)

hFindFile: [in] Search handle returned by a previous call to the FindFirstFile or FindFirstFileEx function.

IpFindFileData: [out] Pointer to the WIN32_FIND_DATA structure that receives information about the found file or subdirectory. The structure can be used in subsequent calls to FindNextFile to indicate from which file to continue the search.

http://msdn2.microsoft.com/en-us/library/aa364418.aspx, http://msdn2.microsoft.com/en-us/library/aa364428.aspx

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Searching for Files (2)

```
source: http://msdn2.microsoft.com/
   #define _WIN32_WINNT 0x0501
                                                                        en-us/library/aa365200.aspx
   #include <windows.h>
  #include <stdio.h>
  #include <streafe h>
  #include <malloc.h>
  #define BUFSIZE MAX_PATH
  int main(int argc, char *argv[]) {
     WIN32_FIND_DATA FindFileData;
     HANDLE hFind = INVALID HANDLE VALUE;
     DWORD dwError; LPSTR DirSpec;
     size t length of arg;
     DirSpec = (LPSTR) malloc (BUFSIZE);
     // Check for command-line parameter; otherwise, print usage.
     if(argc != 2) { printf("Usage: Test <dir>\n"); return 2; }
     // Check that the input is not larger than allowed.
     StringCbLength(argv[1], BUFSIZE, &length_of_arg);
     if (length of arg > (BUFSIZE - 2)) {
        printf("Input directory is too large.\n");
        return 3;
     printf ("Target directory is %s.\n", argv[1]);
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                                                                      9. File Systems (4) - Slide 13
```

Searching for Files (3)

```
// Prepare string for use with FindFile functions. First, copy the
   // string to a buffer, then append '\*' to the directory name.
   StringCbCopyN (DirSpec, BUFSIZE, argv[1], length_of_arg+1);
   StringCbCatN (DirSpec, BUFSIZE, "\\*", 3);
   // Find the first file in the directory.
   hFind = FindFirstFile(DirSpec, &FindFileData);
   if (hFind == INVALID HANDLE VALUE) {
      printf ("Invalid file handle. Error is %u.\n", GetLastError()); return (-1);
      printf ("First file name is %s.\n", FindFileData.cFileName);
      // List all the other files in the directory.
      while (FindNextFile(hFind, &FindFileData) != 0) {
         printf ("Next file name is %s.\n", FindFileData.cFileName);
      dwError = GetLastError();
      FindClose(hFind);
      if (dwError != ERROR_NO_MORE_FILES) {
         printf ("FindNextFile error. Error is %u.\n", dwError); return (-1);
   free(DirSpec);
   return (0);
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                                                                    9. File Systems (4) - Slide 14
```

Windows: stdin, stdout, stderr

- there are three standard file descriptors (as on Linux/Unix)
 - Standard Input (0, stdin)
 - Standard Output (1, stdout)
 - Standard Error Output (2, stderr)
- These can be accessed the same way as on a Linux system:

```
#include <io.h> /* Linux: <unistd.h> */
void main() {
  if (write(1, "Here is some data\n", 18) != 18)
    write(2, "A write error has occurred on file descriptor 1\n",46);
}
```

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File Organization (1)

Criteria for File Organization

- Short access time
 - Needed when accessing a single record
 - Not needed for batch mode
- Ease of update
 - Files on CD-ROM are not updated
- Economy of storage
 - there should be minimum redundancy in the data
 - use redundancy to speed access, e.g. an index
- Simple maintenance
- Reliability

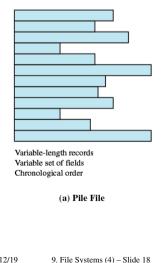
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9. File Systems (4) - Slide 17

File Organization (2)

- The Pile
 - Data are collected in the order they arrive
 - Purpose is to accumulate a mass of data and save it
 - Records may have different fields
 - No structure
 - Record access is by exhaustive search



File Organization (3)

• Example for **Pile**: contents of typical configuration files (here: Wine configuration)

```
[Drive A]
"Type" = "floppy"
                                            "Path"="/media/cdrecorder"
"Path" = "/media/floppy"
                                            "Label"="cdrecorder"
"Label" = "/media/floppy"
                                            "Type"="cdrom"
                                            "FS"="win95"
"Device" = "/dev/fd0"
                                            ; "Device" = " /dev / cdrecorder "
[Drive C]
"Path" = "${HOME}/.wine/fake windows"
                                           [Drive E]
"Type" = "hd"
                                            "Path"="/media/cdrom"
                                            "Label"="cdrom"
"Label" = "fake_windows"
"Filesystem" = "win95"
                                            "Type"="cdrom"
                                            "FS"="win95"
                                            ; "Device"="/dev/cdrom"
```

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9. File Systems (4) - Slide 19

File Organization (4)

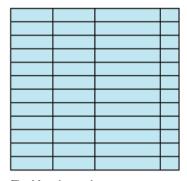
- The Sequential File
 - Fixed format used for records
 - Records are the same length
 - All fields the same (order and length)
 - Field names and lengths are attributes of the file
 - One field is the key field
 - · Uniquely identifies the record
 - · Records are stored in key sequence

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File Organization (5)

- The Sequential File
 - New records are placed in a log file or transaction file
 - Batch update is performed to merge the log file with the master file



Fixed-length records Fixed set of fields in fixed order Sequential order based on key field

(b) Sequential File

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9. File Systems (4) - Slide 21

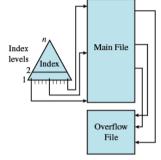
9. File Systems (4) - Slide 22

File Organization (6)

example for sequential file:

File Organization (7)

- Indexed Sequential File
 - Index provides a lookup capability to quickly reach the vicinity of the desired record
 - Contains key field and a pointer to the main file
 - Index is searched to find highest key value that is equal to or precedes the desired key value
 - Search continues in the main file at the location indicated by the pointer



(c) Indexed Sequential File

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9. File Systems (4) - Slide 23

File Organization (8)

- Comparison of sequential and indexed sequential files
 - Example: a file contains 1 million records
 - On average 500,000 accesses are required to find a record in a sequential file
 - If an index contains 1000 entries, it will take an average 500 accesses to find the key, followed by 500 accesses in the main file.
 - Now on average it takes 1000 accesses (that's a factor of 1/500)

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File Organization (9)

- Indexed Sequential File
 - New records are added to an overflow file
 - Record in main file that precedes it is updated to contain a pointer to the new record
 - The overflow is merged with the main file during a batch update
 - Multiple indexes for the same key field can be created to increase efficiency

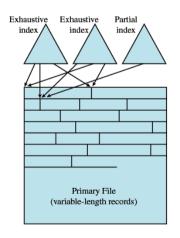
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9. File Systems (4) - Slide 25

File Organization (10)

- Indexed File
 - Uses multiple indexes for different key fields
 - May contain an exhaustive index that contains one entry for every record in the main file
 - May contain a partial index (only entries for records in which this field exists)



(d) Indexed File

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9. File Systems (4) - Slide 26

Secondary Memory Management (1)

- Space must be allocated to files
- Must keep track of the space available for allocation
- Pre-Allocation
 - Need the maximum size for the file at the time of creation
 - Difficult to reliably estimate the maximum potential size of the file
 - Tend to over-estimate file size so as not to run out of space

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9. File Systems (4) - Slide 27

Secondary Memory Management (2)

- Contiguous allocation
 - Single set of blocks is allocated to a file at the time of creation
 - Only a single entry in the file allocation table
 - Starting block and length of the file
- External fragmentation will occur
 - Need to perform compaction

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Secondary Memory Management (3)

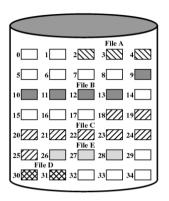




Figure 12.7 Contiguous File Allocation

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9. File Systems (4) - Slide 29

Secondary Memory Management (5)

- Chained allocation
 - Allocation on basis of individual block
 - Each block contains a pointer to the next block in the chain
 - Only single entry in the file allocation table (Starting block and length of file)
 - No external fragmentation
 - Best for sequential files
 - No accommodation of the principle of locality

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9. File Systems (4) - Slide 31

Secondary Memory Management (4)

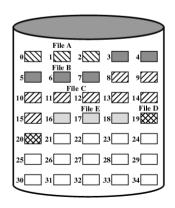




Figure 12.8 Contiguous File Allocation (After Compaction)

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9. File Systems (4) – Slide 30

Secondary Memory Management (6)

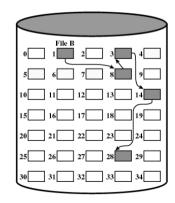




Figure 12.9 Chained Allocation

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Secondary Memory Management (7)

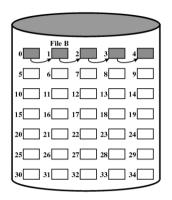




Figure 12.10 Chained Allocation (After Consolidation)

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9. File Systems (4) - Slide 33

9. File Systems (4) - Slide 34

Secondary Memory Management (8)

- Indexed allocation
 - File allocation table contains a separate one-level index for each file
 - The index has one entry for each portion allocated to the file (portions are the smallest units which can be allocated - this size can vary from single blocks to the whole file)
 - The file allocation table contains the block number of the index

Secondary Memory Management (9)

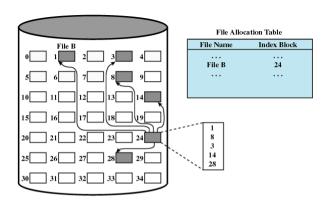


Figure 12.11 Indexed Allocation with Block Portions

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9. File Systems (4) - Slide 35

Secondary Memory Management (10)

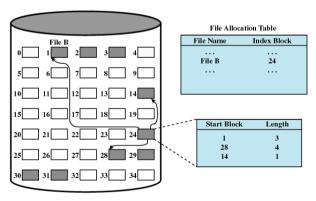


Figure 12.12 Indexed Allocation with Variable-Length Portions

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Secondary Memory Management (11)

Management of free disk space

- manage the disk space which is currently unused
- in addition to FAT (file allocation table) there is a disk allocation table with information about free areas
- bit table: bit vector which each bit representing one disk block:

0 = free block 1 = used block (cf. bit vectors for free main memory management) -> bit table may completely fit in RAM (efficient)

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9. File Systems (4) - Slide 37

Secondary Memory Management (12)

- free block list: blocks are numbered; a list contains the numbers of all free blocks
 - list is too big to be kept in memory completely
 - remedy: treat list as a FIFO queue
 - keep the first few 1000 and the last few 1000 queue entries in memory
 - block allocation: take an entry from the queue's head
 - Block deallocation: append an entry to the queue's tail
 - from time to time synchronize heads and tails with data on disk