

9.7.1 Standard File Systems: Linux Ext2 / Ext3

7 Standard File Systems: Ext2/Ext3 and NTFS

Ext2/Ext3 File System

- Ext3 is today's standard Linux file system (besides ReiserFS)
 - History
 - in the beginning there was Minix...
 - Extended Filesystem replaced Minix
 - Ext2 (2nd Extended Filesystem) succeeded Ext
 - Ext3 (3rd Extended Filesystem) extends Ext2 with Journaling

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Ext2/Ext3: Features (1)

- Features

- variable block sizes (depending on fs creation options, it handles many small or few big files well)
 - fast symbolic links (save link target directly in the inode when the address is short)
 - special attributes (which go beyond the traditional Unix file attributes), e.g. *immutable* flag
 - extended attributes, e.g. ACLs
 - compatibility bitmap for filesystem extensions (*read/write-*, *read-only-* und *incompatible*-Bits)

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Ext2/Ext3: Features (2)

- Features
 - guarantees regular filesystem checks, based on two properties:
 - mount count: mount filesystem only n times without check, then it must be fsck-ed
 - time since last fsck (maximal interval)
 - secure deletion (by overwriting the data blocks)
 - only Ext3: journaling
 - makes long fsck runs unnecessary in normal operation
 - speeds up fsck check in case of errors

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Ext2/Ext3: Design (2)

Super Block	Group Descriptors	Data block Bitmap	inode Bitmap	inode Table	Data blocks
1 block	n block	1 block	1 block	n blocks	n blocks

- **group descriptor:** block group state (among other data: numbers of free blocks and inodes) (variable length)
 - each block group contains descriptors for *all* block groups
- **data block bitmap:** for each data block one bit (free/non-free)

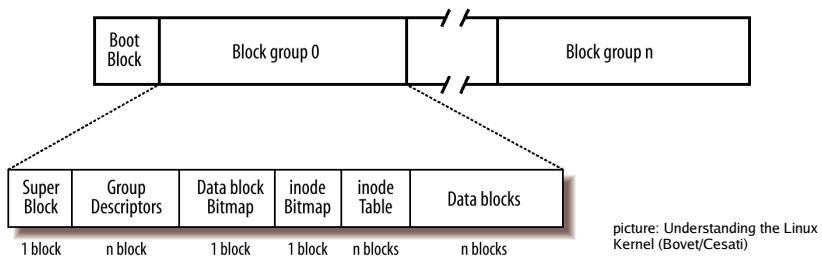
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Ext2/Ext3: Design (1)

- design: partition is split into boot sector and **block groups**
- each block group holds a copy of the **superblock** (with administrative metadata for the whole filesystem)



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Ext2/Ext3: Design (3)

Super Block	Group Descriptors	Data block Bitmap	inode Bitmap	inode Table	Data blocks
1 block	n block	1 block	1 block	n blocks	n blocks

- **inode bitmap:** for each inode one bit: used/unused
- **inode table:** this contains all inodes of the block group (variable length)
- **data blocks:** the real data, files and directories

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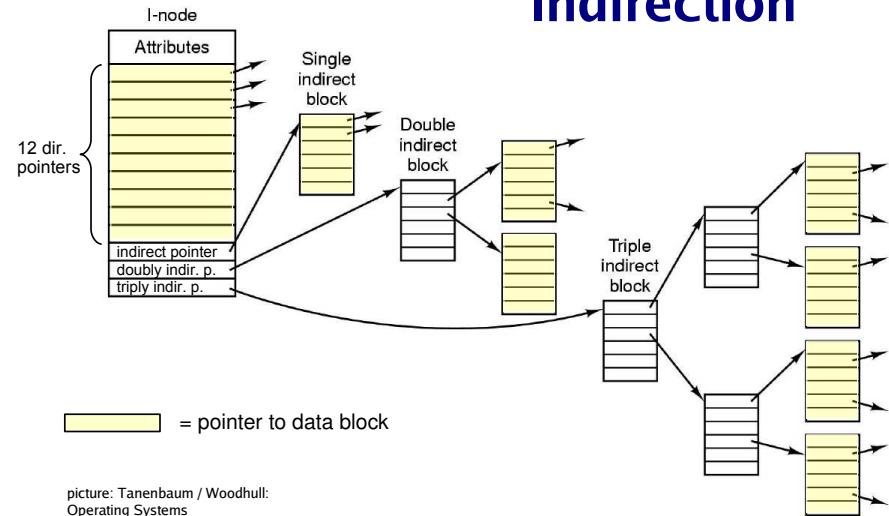
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Ext2/Ext3: Design (4)

- meta information in each block group – idea:
 - if the superblock is destroyed by a failure, there are redundant copies
 - small distance between meta information and data: shorter seek times when accessing the data
- real implementation works differently:
 - kernel keeps a superblock copy in RAM and only distributes it to the other groups when calling *fsck*
 - later Ext2 versions: *Sparse Superblock* option. superblocks only in groups 0, 1, 3^i , 5^i , 7^i ($i > 1$)

File Sizes (2): Multiple Indirection



File Sizes (1)

flexible in two ways:

- when creating the filesystem, several blocksizes are possible:
 - 1024 byte blocks -> max. size 16 GByte
 - 2048 byte blocks -> max. size 256 GByte
 - 4096 byte blocks -> max. size 4096 GByte (4 TB)
- for large files: save block numbers in indirection blocks (up to 3 layers):
 - doubling the block size
= 16 times (2^4) as many block numbers:
factor 2^3 from indirection, factor 2 from block size

File Sizes (3)

Example

length of a (disk) address: 4 Byte

block size BS:	1024 Byte	2048 Byte	4096 Byte
1 block can contain BS/4=N addresses	256	512	1024
12 dir. point.: $12 \cdot BS =$	12 KByte	24 KByte	48 KByte
1x indir. p.: 1 indir. bl.	$256^1 \cdot 1024 = 256 \text{ KB}$	$512^1 \cdot 2048 = 1 \text{ MB}$	$1024^1 \cdot 4096 = 4 \text{ MB}$
2x indir. p.: N indir. bl.	$256^2 \cdot 1024 = 64 \text{ MB}$	$512^2 \cdot 2048 = 512 \text{ MB}$	$1024^2 \cdot 4096 = 4 \text{ GB}$
3x indir. p.: N^2 indir. bl.	$256^3 \cdot 1024 = 16 \text{ GB}$	$512^3 \cdot 2048 = 256 \text{ GB}$	$1024^3 \cdot 4096 = 4 \text{ TB}$

numbers are theoretical; the maximum file size is restricted by other factors as well

Ext2/Ext3 Superblock (1)

- drivers differentiate
 - structure on the **disk** (`ext2_super_block`)
 - structure in **main memory** (`ext2_sb_info`)
- on disk system uses fixed length data structures, e.g.
 - `__u32`: unsigned 32 bit
 - `__s16`: signed 16 bit
- in main memory system uses default datatypes (unsigned long, int etc.)

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Ext2/Ext3 Superblock (2)

- data structures on disk and in memory

Type	Disk data structure	Memory data structure	Caching mode
Superblock	<code>ext2_super_block</code>	<code>ext2_sb_info</code>	Always cached
Group descriptor	<code>ext2_group_desc</code>	<code>ext2_group_desc</code>	Always cached
Block bitmap	Bit array in block	Bit array in buffer	Dynamic
inode bitmap	Bit array in block	Bit array in buffer	Dynamic
inode	<code>ext2_inode</code>	<code>ext2_inode_info</code>	Dynamic
Data block	Array of bytes	VFS buffer	Dynamic
Free inode	<code>ext2_inode</code>	None	Never
Free block	Array of bytes	None	Never

table: Understanding the Linux Kernel (Bovet/Cesati)

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Ext2/Ext3 Superblock (3)

```
struct ext2_super_block {  
    __le32 s_inodes_count;      /* Inodes count */  
    __le32 s_blocks_count;     /* Blocks count */
```

file:
`include/linux/ext2_fs.h`

how many blocks and inodes does the filesystem have?

```
__le32 s_r_blocks_count;   /* Reserved blocks count */  
__le32 s_free_blocks_count; /* Free blocks count */  
__le32 s_free_inodes_count; /* Free inodes count */
```

free inodes and blocks

```
__le32 s_first_data_block; /* First Data Block */  
__le32 s_log_block_size;  /* Block size */
```

logarithmic block size:

$1024 \cdot 2^i$, $i=0,1,2$ (1024, 2048, 4096)

```
__le32 s_log_frag_size; /* Fragment size */  
__le32 s_blocks_per_group; /* # Blocks per group */  
__le32 s frags_per_group; /* # Fragments per group */  
__le32 s_inodes_per_group; /* # Inodes per group */
```

inodes / blocks in a block group

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Ext2/Ext3 Superblock (4)

```
__le32 s_mtime;           /* Mount time */  
__le32 s_wtime;          /* Write time */  
__le16 s_mnt_count;       /* Mount count */  
__le16 s_max_mnt_count;  /* Maximal mount count */
```

mount count is increased each time the fs is mounted; when a maximum value is reached, it is time for an `fsck`

```
/* Maximal mount counts between two filesystem checks */  
#define EXT2_DFL_MAX_MNT_COUNT      20 /* Allow 20 mounts */
```

```
__le16 s_magic;           /* Magic signature */
```

```
/* The second extended file system magic number */
```

```
#define EXT2_SUPER_MAGIC 0xEF53
```

```
__le16 s_state;           /* File system state */
```

```
/*
```

```
* File system states
```

```
*/
```

```
#define EXT2_VALID_FS        0x0001 /* Unmounted cleanly */
```

```
#define EXT2_ERROR_FS        0x0002 /* Errors detected */
```

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Ext2/Ext3 Superblock (5)

```

__le16 s_errors;           /* Behaviour when detecting errors */

/*Behaviour when detecting errors */
#define EXT2_ERRORS_CONTINUE    1 /* Continue execution */
#define EXT2_ERRORS_RO          2 /* Remount fs read-only */
#define EXT2_ERRORS_PANIC       3 /* Panic */
#define EXT2_ERRORS_DEFAULT     EXT2_ERRORS_CONTINUE

__le16 s_minor_rev_level; /* minor revision level */
__le32 s_lastcheck;        /* time of last check */
__le32 s_checkinterval;    /* max. time between checks */

lastcheck and checkinterval: besides having a mount count there
are also automatic fsck checks after a given time interval passed

__le32 s_creator_os;       /* OS */
__le32 s_rev_level;        /* Revision level */
__le16 s_def_resuid;       /* Default uid for reserved blocks */
__le16 s_def_resgid;       /* Default gid for reserved blocks */

„reserved space“ (5 %) for a privileged user.
default: resuid = resgid = 0 (root/root)

#define EXT2_DEF_RESUID      0
#define EXT2_DEF_RESID       0

```

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Ext2/Ext3 Superblock (6)

```

__le32 s_first_ino;          /* First non-reserved inode */
__le16 s_inode_size;         /* size of inode structure */
__le16 s_block_group_nr;     /* block group # of this superblock */
__le32 s_feature_compat;     /* compatible feature set */
__le32 s_feature_incompat;   /* incompatible feature set */
__le32 s_feature_ro_compat;  /* readonly-compatible feature set */

```

three feature categories that the kernel uses for deciding, whether and how it will mount this filesystem:

- compatible: mount filesystem r/w, even if unsupported
- ro-comp.: mount filesystem r/o if unsupported
- incompatible: do not mount filesystem at all if unsupported

`s_feature_compat`, `s_feature_ro_compat`, `s_feature_incompat` are bitmaps which describe the features active *in this filesystem*, e.g.

```

#define EXT2_FEATURE_COMPAT_DIR_PREALLOC 0x0001
#define EXT2_FEATURE_COMPAT_IMAGIC_INODES 0x0002
#define EXT3_FEATURE_COMPAT_HAS_JOURNAL 0x0004
...
#define EXT2_FEATURE_INCOMPAT_COMPRESSION 0x0001
#define EXT2_FEATURE_INCOMPAT_FILETYPE   0x0002
#define EXT3_FEATURE_INCOMPAT_RECOVER    0x0004
...

```

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Ext2/Ext3 Superblock (7)

```

__u8   s_uuid[16];           /* 128-bit uuid for volume */
char  s_volume_name[16];     /* volume name */
char  s_last_mounted[64];    /* directory where last mounted */
__le32 s_algorithm_usage_bitmap; /* For compression */

/*
 * Performance hints. Directory preallocation should only
 * happen if the EXT2_COMPAT_PREALLOC flag is on.
 */
__u8   s_prealloc_blocks;    /* Nr of blocks to try to preallocate */
__u8   s_prealloc_dir_blocks; /* Nr to preallocate for dirs */

```

Pre-allocation: when allocating a block, Ext2 will not only allocate the requested block for the file, but will reserve several blocks in advance: this reduces potential fragmentation and speeds up further block requests for the same file which follow briefly after the current request (e.g. when sequentially writing to a file)

__u16 s_padding1;

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Ext2/Ext3 Superblock (8)

```

/*
 * Journaling support valid if EXT3_FEATURE_COMPAT_HAS_JOURNAL set.
 */
__u8   s_journal_uuid[16]; /* uuid of journal superblock */
__u32  s_journal_inum;    /* inode number of journal file */
__u32  s_journal_dev;     /* device number of journal file */

```

the journal can reside on a different filesystem; this not only a file inode, but also a device number

```

__u32  s_last_orphan;        /* start of list of inodes to delete */
__u32  s_hash_seed[4];       /* HTREE hash seed */
__u8   s_def_hash_version;   /* Default hash version to use */
__u8   s_reserved_char_pad;
__u16  s_reserved_word_pad;
__le32 s_default_mount_opts;
__le32 s_first_meta_bg;      /* First metablock block group */
__u32  s_reserved[190];      /* Padding to the end of the block */
};

```

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Ext2 Superblock in Memory (1)

```
/* second extended-fs super-block data in memory */
struct ext2_sb_info {
    unsigned long s_frag_size; /* Size of a fragment in bytes */
    unsigned long s frags_per_block; /* Number of fragments per block */
    unsigned long s_inodes_per_block; /* Number of inodes per block */
    unsigned long s frags_per_group; /* Number of fragments in a group */
    unsigned long s_blocks_per_group; /* Number of blocks in a group */
    unsigned long s_inodes_per_group; /* Number of inodes in a group */
    unsigned long s_itb_per_group; /* Number of inode table blocks per group */
    unsigned long s_gdb_count; /* Number of group descriptor blocks */
    unsigned long s_desc_per_block; /* Number of group descriptors per block */
    unsigned long s_groups_count; /* Number of groups in the fs */
    struct buffer_head * s_sbh; /* Buffer containing the super block */
    struct ext2_super_block * s_es; /* Pointer to the super block in the buffer */
    struct buffer_head ** s_group_desc;
    unsigned long s_mount_opt;

#define EXT2_MOUNT_CHECK 0x0001 /* Do mount-time checks */
#define EXT2_MOUNT_OLDALLOC 0x0002 /* Don't use the new Orliv allocator */
#define EXT2_MOUNT_GRPID 0x0004 /* Create files with directory's group */
#define EXT2_MOUNT_DEBUG 0x0008 /* Some debugging messages */
#define EXT2_MOUNT_ERRORS_CONT 0x0010 /* Continue on errors */
#define EXT2_MOUNT_ERRORS_RO 0x0020 /* Remount fs ro on errors */
#define EXT2_MOUNT_ERRORS_PANIC 0x0040 /* Panic on errors */
#define EXT2_MOUNT_MINIX_DF 0x0080 /* Mimics the Minix statfs */
#define EXT2_MOUNT_NOBH 0x0100 /* No buffer heads */
#define EXT2_MOUNT_NO_UID32 0x0200 /* Disable 32-bit UIDs */
#define EXT2_MOUNT_XATTR_USER 0x4000 /* Extended user attributes */
#define EXT2_MOUNT_POSIX_ACL 0x8000 /* POSIX Access Control Lists */
}
```

file:
include/linux/
ext2_fs_sb.h

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Ext2 Inode (1)

```
/* Structure of an inode on the disk */
struct ext2_inode {
    __le16 i_mode;
    __le16 i_uid;
    __le32 i_size;
    __le32 i_atime;
    __le32 i_ctime;
    __le32 i_mtime;
    __le32 i_dtime;
    __le16 i_gid;
    __le16 i_links_count;
    __le32 i_blocks;
    __le32 i_flags;
}

Flags: EXT2_SECRM_FL 0x00000001 secure deletion
EXT2_UNRM_FL 0x00000002 record for undelete
EXT2_COMPR_FL 0x00000004 compressed file
EXT2_SYNC_FL 0x00000008 synchronous updates
EXT2_IMMUTABLE_FL 0x00000010 immutable file
EXT2_APPEND_FL 0x00000020 append only
EXT2_NODUMP_FL 0x00000040 do not dump/delete file
EXT2_NOATIME_FL 0x00000080 do not update .i_atime
EXT2_DIRTY_FL 0x00000100 dirty (file is in use?)
EXT2_COMPRFL_FL 0x00000200 compressed blocks
EXT2_NOCOMPR_FL 0x00000400 access raw compressed data
EXT2_ECOMPR_FL 0x00000800 compression error
EXT2_BTREE_FL 0x00010000 b-tree format directory
EXT2_INDEX_FL 0x00010000 Hash indexed directory
EXT3_JOURNAL_DATA_FL 0x00040000 journal file data
EXT2_RESERVED_FL 0x80000000 reserved for ext2 implementation
```

file:
include/linux/ext2_fs.h

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Ext2 Superblock in Memory (2)

```
uid_t s_resuid;
gid_t s_resgid;
unsigned short s_mount_state;
unsigned short s_pad;
int s_addr_per_block_bits;
int s_desc_per_block_bits;
int s_inode_size;
int s_first_ino;
spinlock_t s_next_gen_lock;
u32 s_next_generation;
unsigned long s_dir_count;
u8 *s_debts;
struct percpu_counter s_freeblocks_counter;
struct percpu_counter s_freeinodes_counter;
struct percpu_counter s_dirs_counter;
struct blockgroup_lock s_blockgroup_lock;
};

}
```

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Ext2 Inode (2)

```
union {
    struct {
        __le32 l_i_reserved1;
    } linux1;
    [...]
} osd1; /* OS dependent 1 */
__le32 i_block[EXT2_N_BLOCKS]; /* Pointers to blocks */
__le32 i_generation; /* File version (for NFS) */
__le32 i_file_acl; /* File ACL */
__le32 i_dir_acl; /* Directory ACL */

#define i_size_high i_dir_acl

for large files with a 64 bit length, i_dir_acl (which is otherwise unused for files) is used as upper 32 bits, i_size lower 32 bits

__le32 i_faddr; /* Fragment address */
union {
    struct {
        __u8 l_i_frag; /* Fragment number */
        __u8 l_i_fsize; /* Fragment size */
        __u16 i_padi;
        __le16 l_i_uid_high; /* these 2 fields */
        __le16 l_i_gid_high; /* were reserved2[0] */
        __u32 l_i_reserved2;
    } linux2;
    [...]
} osd2; /* OS dependent 2 */
};
```

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Extra Flags

- Ext2/Ext3 extra flags (immutable, append-only etc.) can be modified with **chattr** (cf. 9. File Systems (2)/Folie 21)

```
NAME
    chattr - change file attributes on a Linux second extended file system

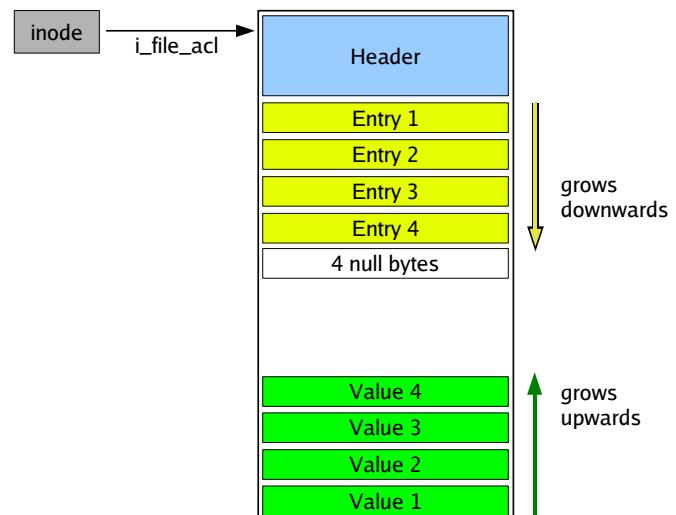
SYNOPSIS
    chattr [ -RV ] [ -v version ] [ mode ] files...

DESCRIPTION
    chattr changes the file attributes on a Linux second extended file system.

    The format of a symbolic mode is +-=[ASacDdIijsTtu].

    The operator `+' causes the selected attributes to be added to the existing
    attributes of the files; `-' causes them to be removed; and `=' causes them
    to be the only attributes that the files have.
```

Extended Attributes (2)



Extended Attributes (1)

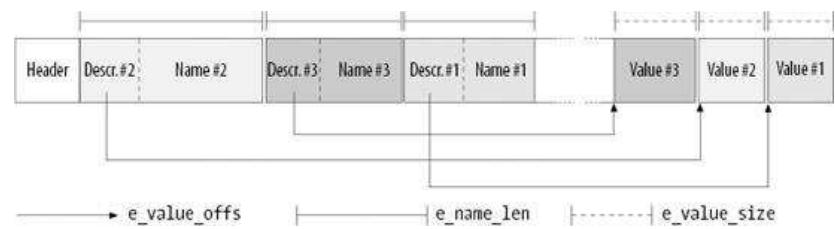
- inode size: 128 bytes
 - no space for extended attributes
 - increasing inodes to 256 bytes is inefficient
- solution: separate block for extended attributes (`ext2_xattr_entry`)
 - xattr header:

```
file: fs/ext2/xattr.h
struct ext2_xattr_header {
    __le32 h_magic;           /* magic number for identification */
    __le32 h_refcount;        /* reference count */
    __le32 h_blocks;          /* number of disk blocks used */
    __le32 h_hash;            /* hash value of all attributes */
    __u32 h_reserved[4];      /* zero right now */
};
```

Extended Attributes (3)

xattr entry:

```
file: fs/ext2/xattr.h
struct ext2_xattr_entry {
    __u8 e_name_len;          /* length of name */
    __u8 e_name_index;        /* attribute name index */
    __le16 e_value_offs;       /* offset in disk block of value */
    __le32 e_value_block;     /* disk block attribute is stored on (n/i) */
    __le32 e_value_size;      /* size of attribute value */
    __le32 e_hash;            /* hash value of name and value */
    char e_name[0];           /* attribute name */
};
```



Extended Attributes (4)

- modify and query with **setfattr**, **getfattr**, **attr**:

```
amd64:/home/esser # setfattr -n user.foo -v betriebssysteme test.txt
amd64:/home/esser # getfattr -d test.txt
# file: test.txt
user.foo="betriebssysteme"

amd64:/home/esser # attr -g user.foo test.txt
Attribute "user.foo" had a 5 byte value for test.txt:
hallo
```

Three Kinds of Attributes

don't confuse:

- standard Unix file attributes
 - UID, GID
 - standard access rights rwx,
 - access times, ...
- extra flags
 - immutable, compressed, secure deletion, ...
- extended attributes
 - arbitrary, freely-definable attributes (incl. ACLs)